

GLOBAL OPTIMIZATION WITH GAMS/BARON



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MIXED-INTEGER NONLINEAR PROGRAMMING

$$\begin{aligned} \text{(P)} \quad & \min f(x, y) \\ \text{s.t.} \quad & g(x, y) \leq 0 \\ & x \in \mathbb{R}^n \\ & y \in \mathbb{Z}^p \end{aligned}$$

Objective Function

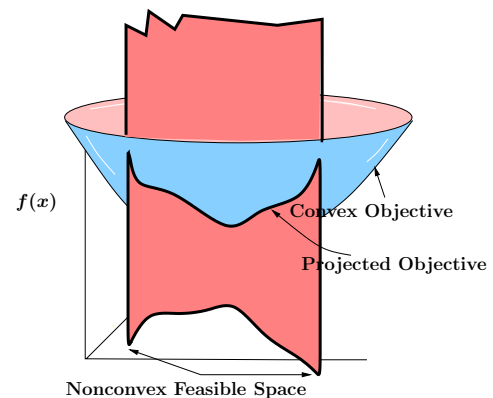
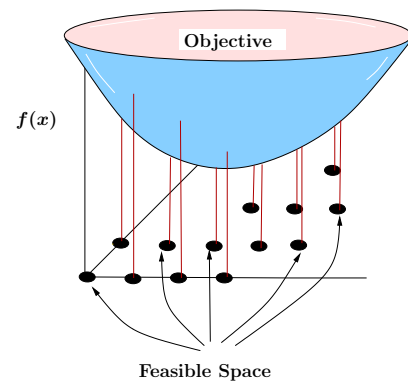
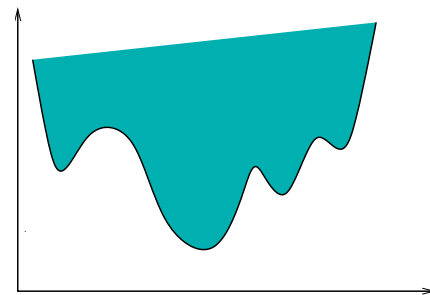
Constraints

Continuous Variables

Integrality Restrictions

Challenges:

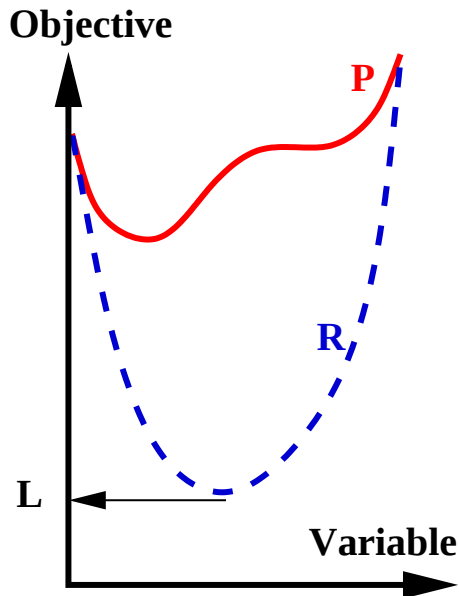
- Multimodal Objective
- Integrality
- Nonconvex Constraints



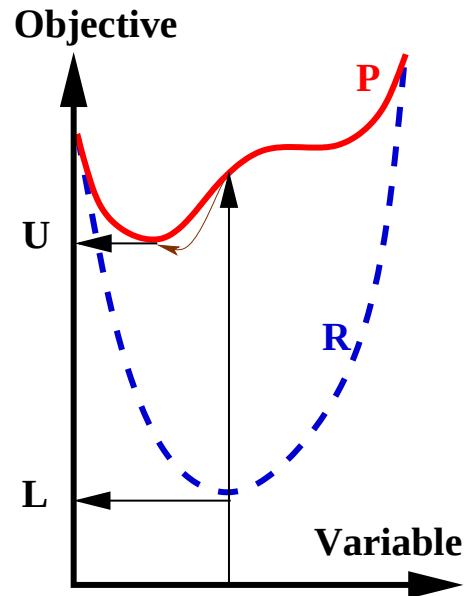
MINLP ALGORITHMS

- **Branch-and-Bound**
 - Bound problem over successively refined partitions
 - » Falk and Soland, 1969
 - » McCormick, 1976
- **Convexification**
 - Outer-approximate with increasingly tighter convex programs
 - Tuy, 1964
 - Serali and Adams, 1994
- **Decomposition**
 - Project out some variables by solving subproblem
 - » Duran and Grossmann, 1986
 - » Visweswaran and Floudas, 1990
- **Our approach**
 - Branch-and-Reduce
 - » Ryoo and Sahinidis, 1995, 1996
 - » Sheckman and Sahinidis, 1998
 - Constraint Propagation & Duality-Based Reduction
 - » Ryoo and Sahinidis, 1995, 1996
 - » Tawarmalani and Sahinidis, 2002
 - Convexification
 - » Tawarmalani and Sahinidis, 2001, 2002
- **Tawarmalani, M. and N. V. Sahinidis, *Convexification and Global Optimization in Continuous and Mixed-Integer Nonlinear Programming*, Kluwer Academic Publishers, Nov. 2002.**

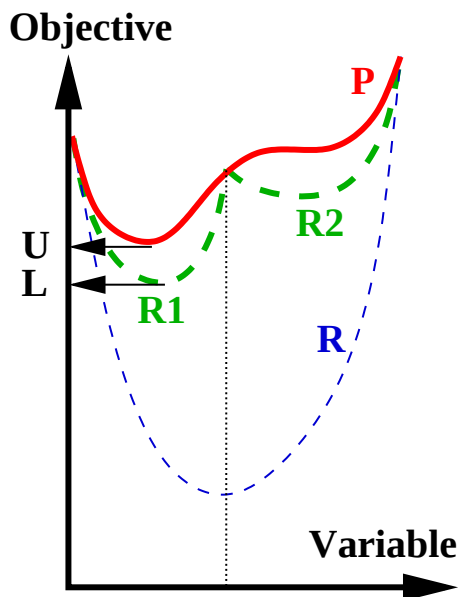
BRANCH-AND-BOUND



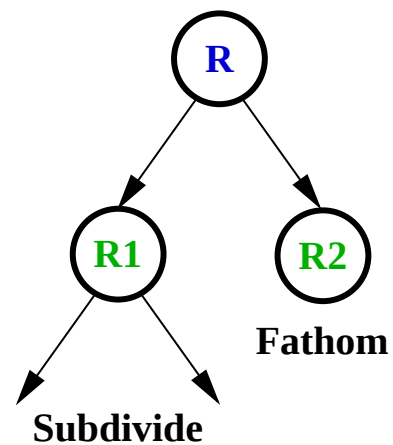
a. Lower Bounding



b. Upper Bounding



c. Domain Subdivision



d. Search Tree

FACTORABLE FUNCTIONS

(McCormick, 1976)

Definition: Factorable functions are recursive compositions of sums and products of functions of single variables.

Example: $f(x, y, z, w) = \sqrt{\exp(xy + z \ln w) z^3}$

The diagram illustrates the recursive construction of the function f from its variables x, y, z, w using intermediate variables x_1 through x_7 . The expression is shown with curly braces and labels indicating the sequence of operations:

- x_1 is assigned to xy .
- x_2 is assigned to $\ln w$.
- x_3 is assigned to zx_2 .
- x_4 is assigned to $x_1 + x_3$.
- x_5 is assigned to $\exp(x_4)$.
- x_6 is assigned to z^3 .
- x_7 is assigned to $x_5 x_6$.
- The final function f is defined as $\sqrt{x_7}$.

$$x_1 = xy$$

$$x_2 = \ln(w)$$

$$x_3 = zx_2$$

$$x_4 = x_1 + x_3$$

$$x_5 = \exp(x_4)$$

$$x_6 = z^3$$

$$x_7 = x_5 x_6$$

$$f = \sqrt{x_7}$$

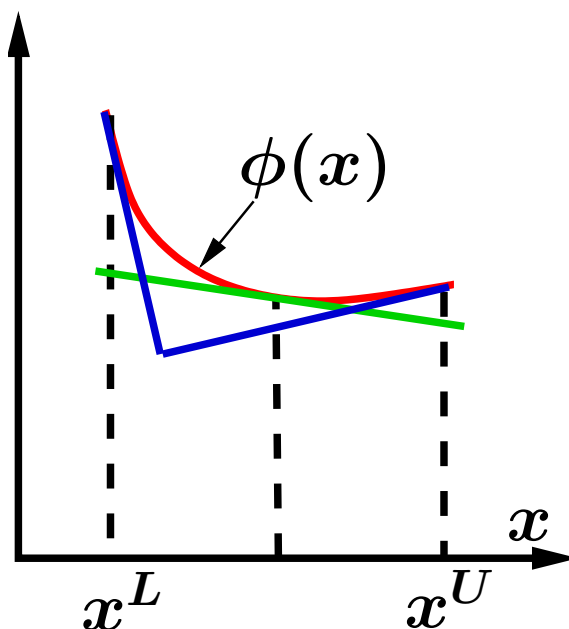
OUTER APPROXIMATION

Motivation:

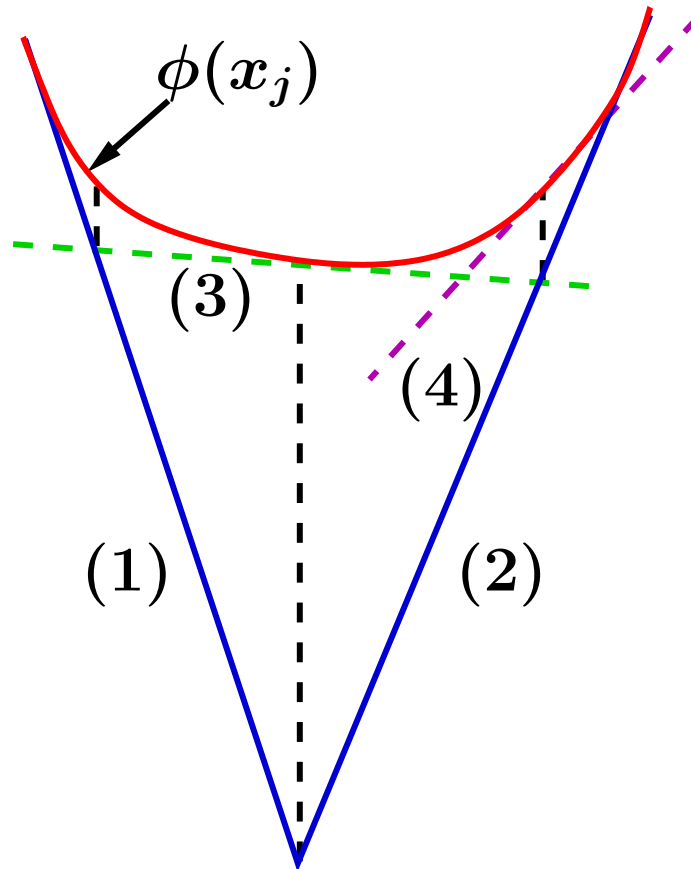
- Convex NLP solvers are not as robust as LP solvers
- Linear programs can be solved efficiently

Outer-Approximation:

Convex Functions are underestimated by tangent lines



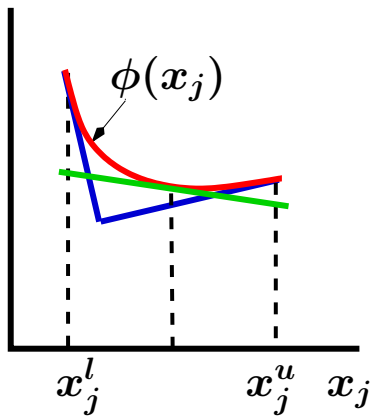
THE SANDWICH ALGORITHM



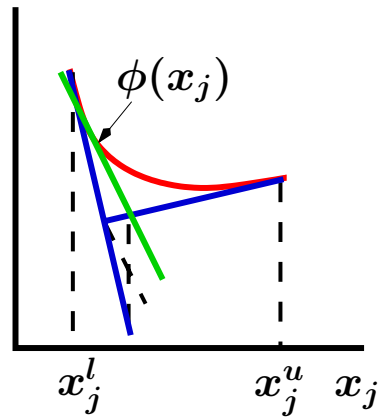
An adaptive strategy

- Assume an initial outer-approximation
- Find point maximizing an error measure
- Construct underestimator at located point

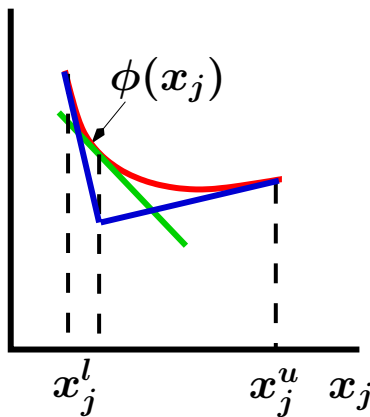
TANGENT LOCATION RULES



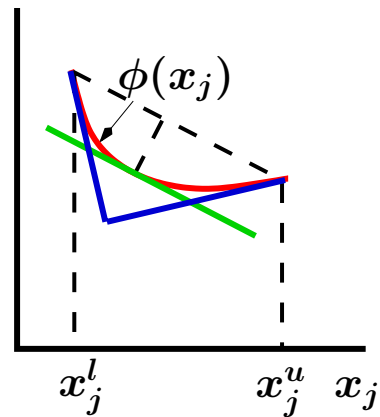
Interval bisection



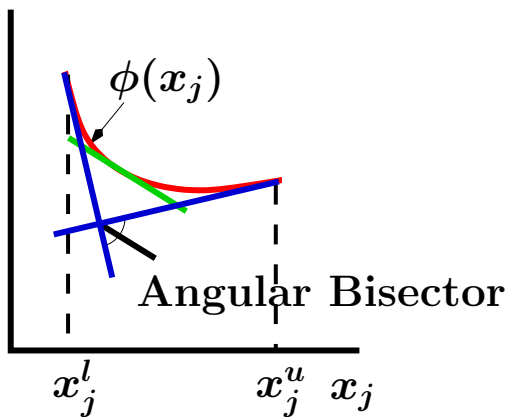
Slope Bisection



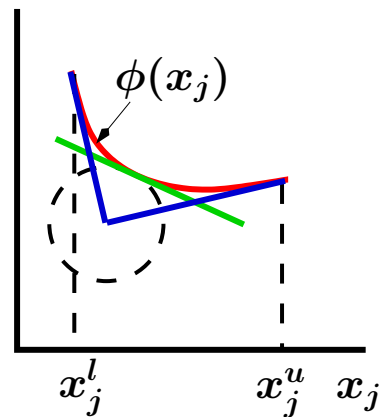
Maximum error rule



Chord rule



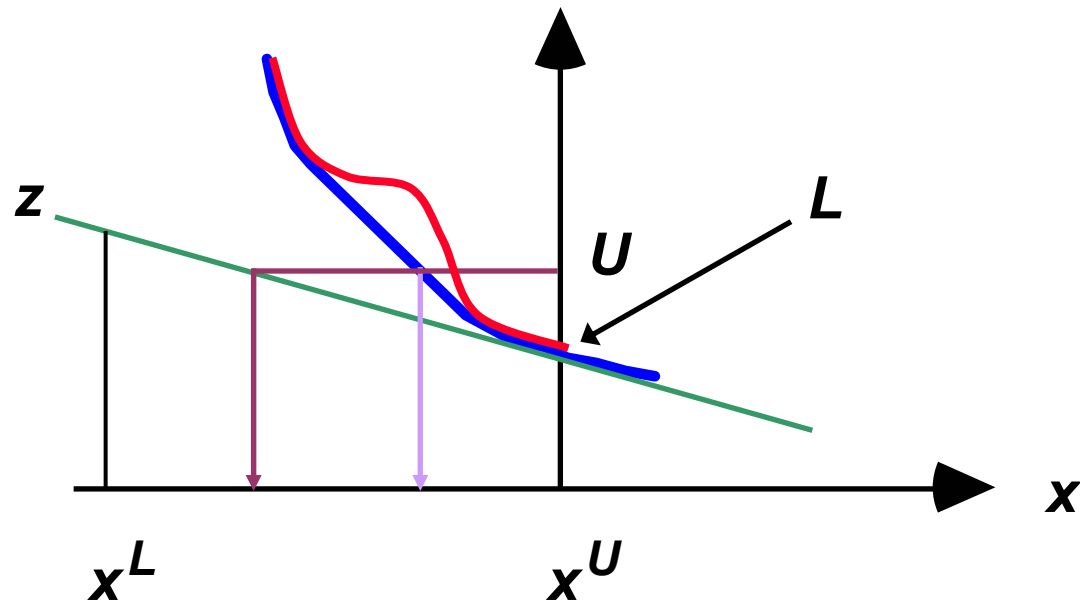
Angle bisection



Maximum projective error

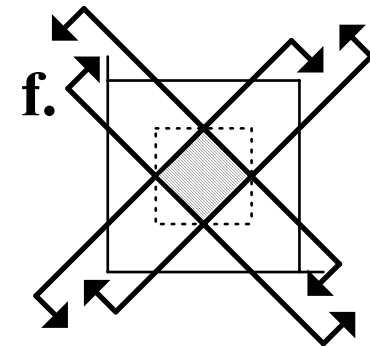
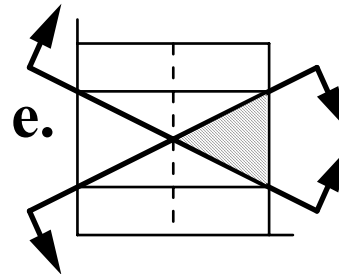
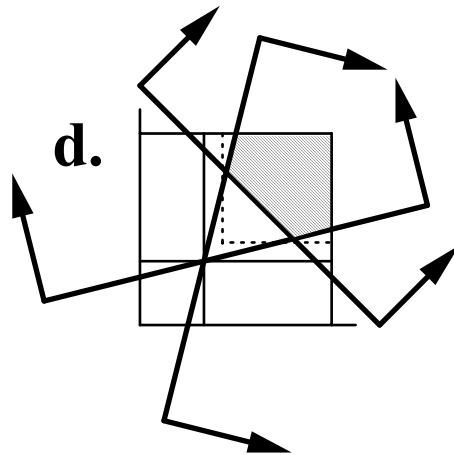
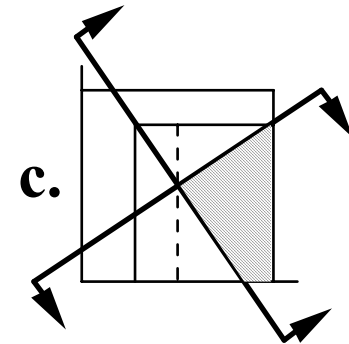
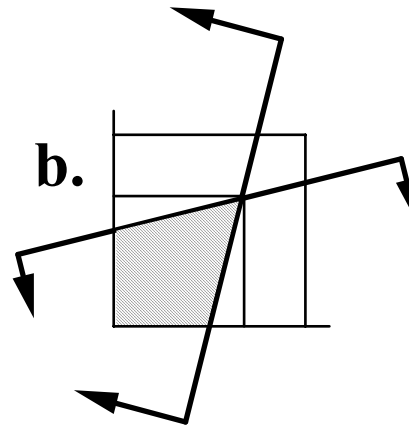
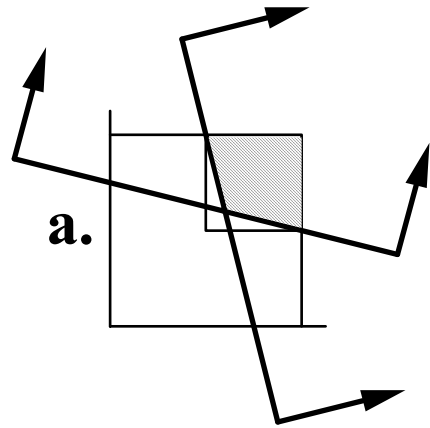
MARGINALS-BASED RANGE REDUCTION

Relaxed Value Function



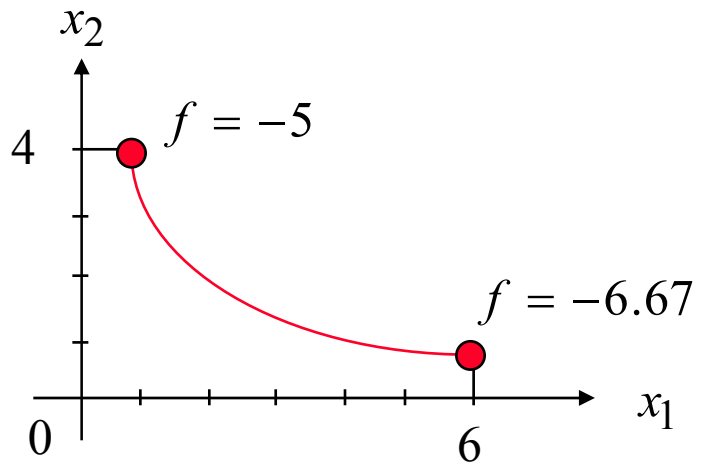
If a variable goes to its upper bound at the relaxed problem solution, this variable's lower bound can be improved

“POOR MAN’S LPs AND NLPs”



ILLUSTRATIVE EXAMPLE

$$\begin{array}{ll} \min & f = -x_1 - x_2 \\ \text{s.t.} & \left. \begin{array}{l} x_1 x_2 \leq 4 \\ 0 \leq x_1 \leq 6 \\ 0 \leq x_2 \leq 4 \end{array} \right\} (P) \end{array}$$



Reformulation :

$$\begin{array}{ll} \min & f = -x_1 - x_2 \\ \text{s.t.} & \left. \begin{array}{l} x_3^2 - x_1^2 - x_2^2 \leq 8 \\ x_3 = x_1 + x_2 \\ 0 \leq x_1 \leq 6 \\ 0 \leq x_2 \leq 4 \\ 0 \leq x_3 \leq 10 \end{array} \right\} \end{array}$$

Relaxation :

$$\begin{array}{ll} \min & -x_1 - x_2 \\ \text{s.t.} & \left. \begin{array}{l} x_3^2 - 6x_1 - 4x_2 \leq 8 \\ x_3 = x_1 + x_2 \\ 0 \leq x_1 \leq 6 \\ 0 \leq x_2 \leq 4 \\ 0 \leq x_3 \leq 10 \end{array} \right\} (R) \end{array}$$

Solution of R : $x_1 = 6, x_2 = 0.89, L = -6.89, \lambda_1 = 0.2$

Local Solution of P with MINOS : $U = -6.67$

Range reduction : $x_1^L \leftarrow x_1^U - (U - L) / \lambda_1 = 4.86$

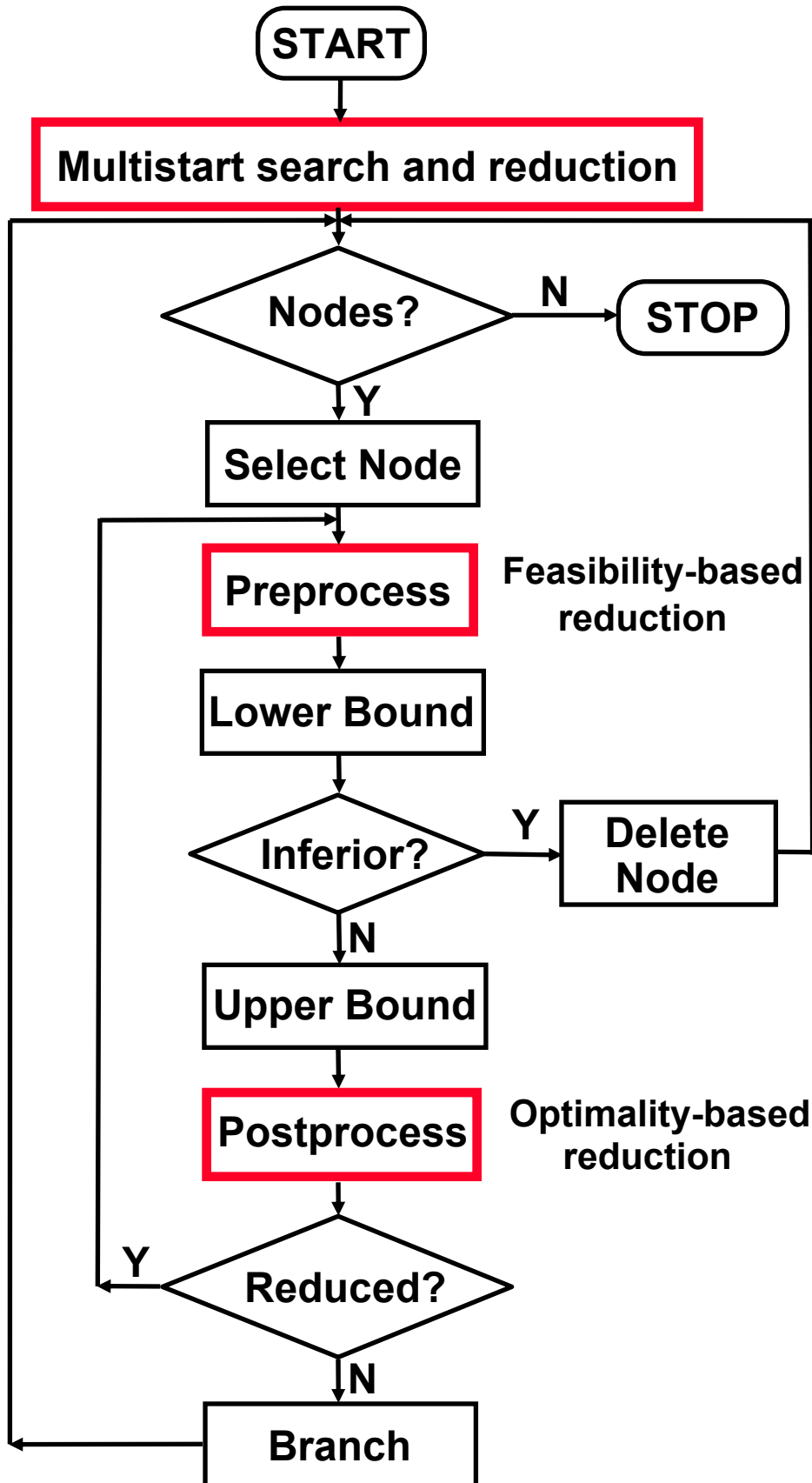
Probing (Solve R with $x_2 \leq 0$) : $L = -6, \lambda_2 = 1 \Rightarrow x_2^L = 0.67$

Update R with $x_1 \geq 4.86, x_2 \geq 0.67$

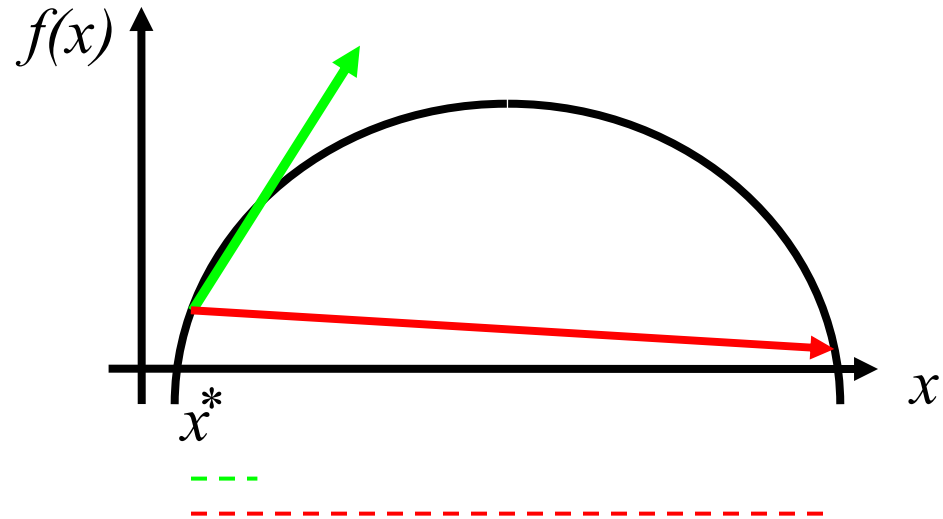
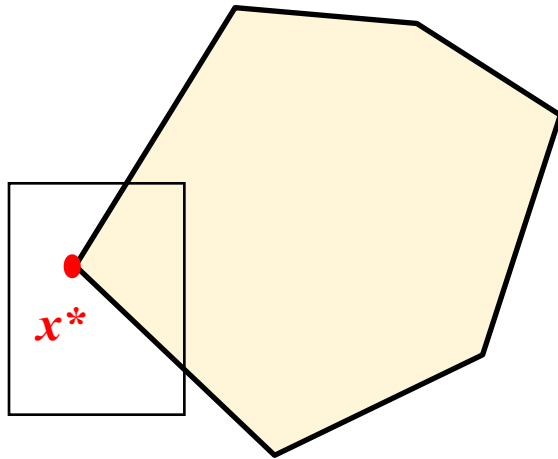
Solution is : $L = -6.67$

$\therefore$ Proof of globality with NO Branching!!

Branch-and-REDUCE



FINITE BRANCHING RULE

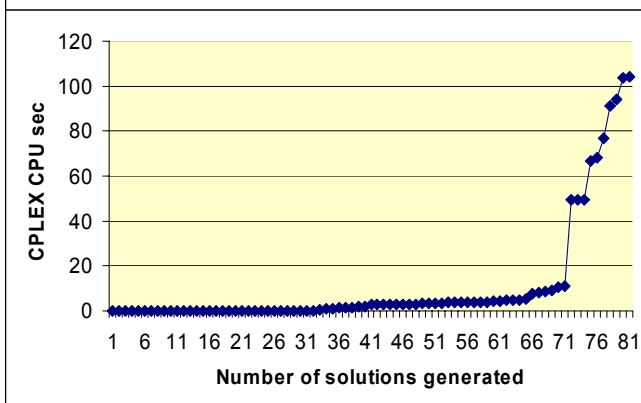
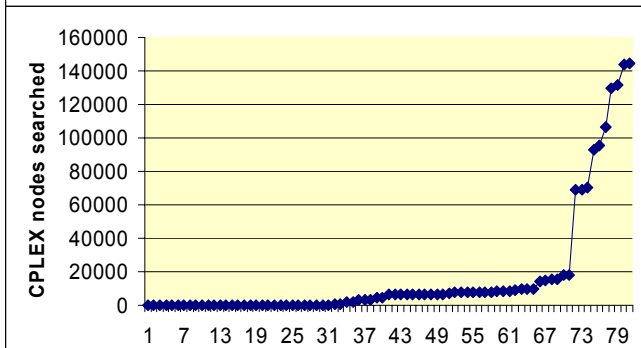
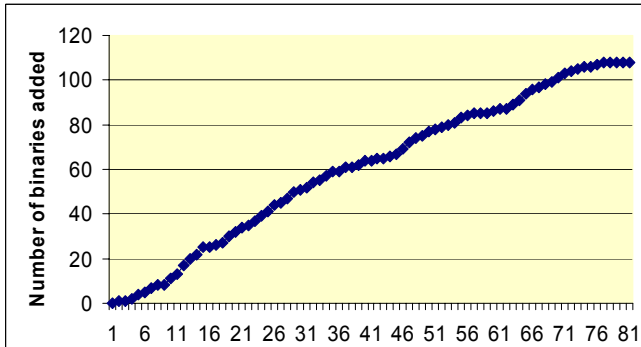


- **Variable selection:**
 - Typically, select variable with largest underestimating gap
 - Occasionally, select variable corresponding to largest edge
- **Locating partition:**
 - Typically, at the midpoint (exhaustiveness)
 - When possible, at the best currently known solution
 - » Makes underestimators exact at the candidate solutions
- **Finite isolation of global optimum**
- **Finite termination in many cases**

FINDING THE K -BEST OR ALL FEASIBLE SOLUTIONS

Typically found through repetitive applications of branch-and-bound and generation of “integer cuts”

$$\begin{aligned} \min \quad & \sum_{i=1}^4 10^{4-i} x_i \\ \text{s.t.} \quad & 2 \leq x_i \leq 4, \quad i = 1, \dots, 4 \\ & x \text{ integer} \end{aligned}$$

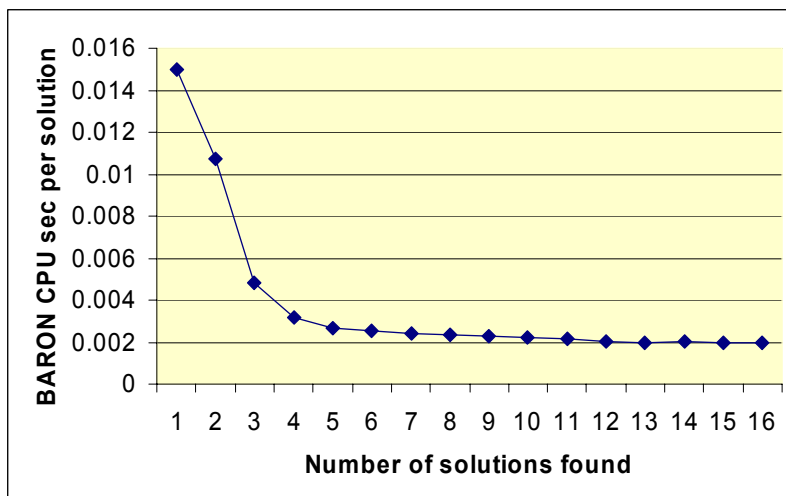
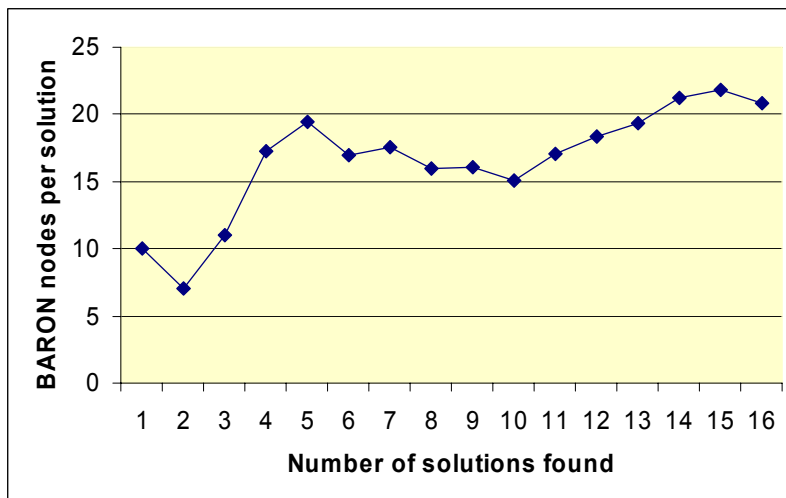


BARON finds all solutions:

- No integer cuts
- Fathom nodes that are infeasible or points
- Single search tree
- 511 nodes; 0.56 seconds
- Applicable to discrete and continuous spaces

FINDING ALL or the K-BEST SOLUTIONS for CONTINUOUS PROBLEMS

- Boon problem: 8 solutions (3.1 sec)
- Robot problem: 16 solutions (0.03 sec)



Branch-And-Reduce Optimization Navigator

Components

- Modeling language
- Preprocessor
- Data organizer
- I/O handler
- Range reduction
- Solver links
- Interval arithmetic
- Sparse matrix routines
- Automatic differentiator
- IEEE exception handler
- Debugging facilities

Capabilities

- Core module
 - Application-independent
 - Expandable
- Fully automated MINLP solver
- Application modules
 - Multiplicative programs
 - Indefinite QPs
 - Fixed-charge programs
 - Mixed-integer SDPs
 - ...
- Solve relaxations using
 - CPLEX, MINOS, SNOPT, OSL, SDPA, ...

- First on the Internet in March 1995
- On-line solver between October 1999 and May 2003
 - Solved eight problems a day
- GAMS link since November 2000

GUPTA-RAVINDRAN MINLPs

Problem		Obj.	T_{tot}	N_{tot}	N_{mem}
1		12.47	0.11	22	4
2	*	5.96	0.03	7	4
3		16.00	0.03	3	2
4		0.72	0.01	1	1
5		5.47	4.48	232	22
6		1.77	0.06	11	5
7		4.00	0.03	3	2
8		23.45	0.40	7	2
9		-43.13	0.58	37	7
10		-310.80	0.06	12	4
11		-431.00	0.12	34	8
12		-481.20	0.29	67	12

Problem		Obj.	T_{tot}	N_{tot}	N_{mem}
13		-585.20	1.13	197	28
14	*	-40358.20	0.05	7	4
15		1.00	0.05	11	3
16		0.70	0.05	23	12
17		-1100.40	42.2	3489	399
18		-778.40	8.85	993	121
19		-1098.40	133	6814	833
20	*	230.92	6.58	143	18
21	*	-5.68	0.21	54	5
22		6.06	2.36	171	39
23		-1125.20	1152	39918	4678
24		-1033.20	4404	124282	15652

* Indicates that a better solution was found than reported in Gupta and Ravindran, Man. Sci., 1985.

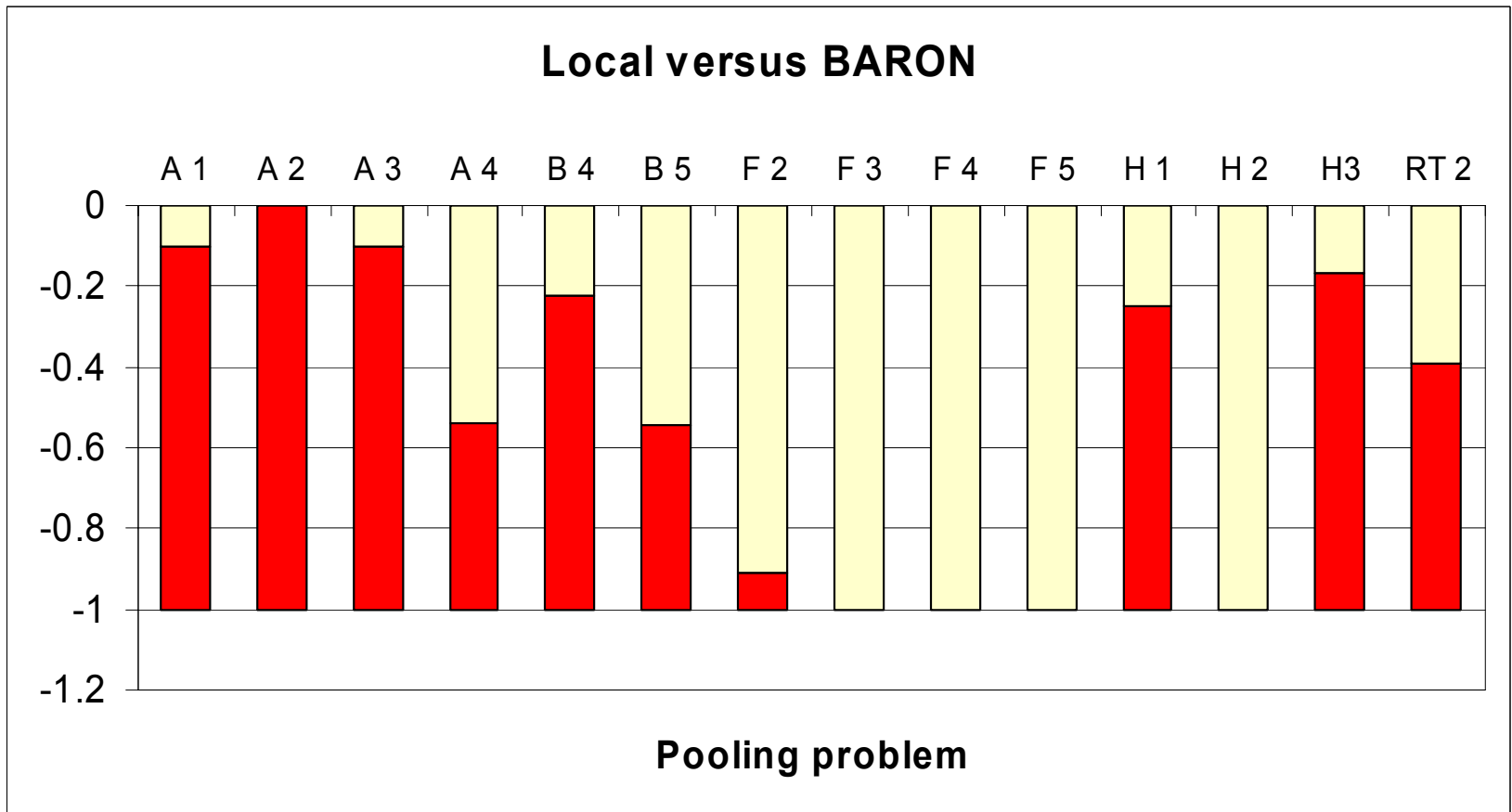
POOLING PROBLEMS

Algorithm	Foulds '92		Ben-Tal '94		GOP '96		BARON '99		BARON '01	
Computer*	CDC 4340				HP9000/730		RS6000/43P		RS6000/43P	
Linpac	> 3.5				49		59.9		59.9	
Tolerance*					**		10^{-6}		10^{-6}	
Problem	N_{tot}	T_{tot}	N_{tot}	T_{tot}	N_{tot}	T_{tot}	N_{tot}	T_{tot}	N_{tot}	T_{tot}
Haverly 1	5	0.7	3	-	12	0.22	3	0.09	1	0.09
Haverly 2			3	-	12	0.21	9	0.09	1	0.13
Haverly 3			3	-	14	0.26	5	0.13	1	0.07
Foulds 2	9	3.0					1	0.10	1	0.04
Foulds 3	1	10.5					1	2.33	1	1.70
Foulds 4	25	125.0					1	2.59	1	0.38
Foulds 5	125	163.6					1	0.86	1	0.10
Ben-Tal 4			25	-	7	0.95	3	0.11	1	0.13
Ben-Tal 5			283	-	41	5.80	1	1.12	1	1.22
Adhya 1							6174	425	15	4.00
Adhya 2							10743	1115	19	4.48
Adhya 3							79944	19314	5	3.16
Adhya 4							1980	182	1	0.97

* Blank indicates problem not reported or not solved

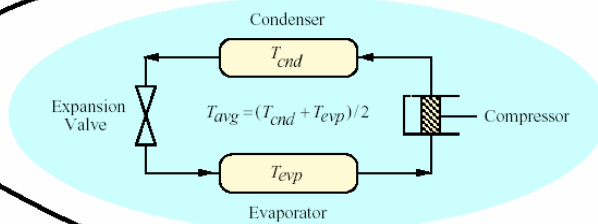
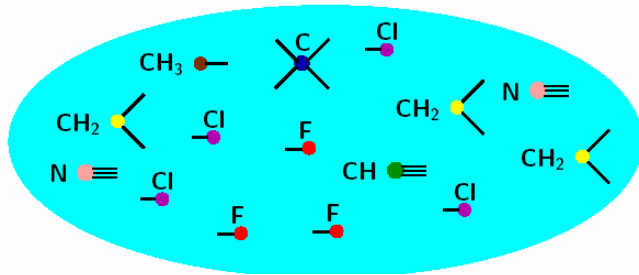
** 0.05% for Haverly 1, 2, 3, 0.05% for Ben-Tal 4 and 1% for Ben-Tal 5

LOCAL vs. GLOBAL SEARCH

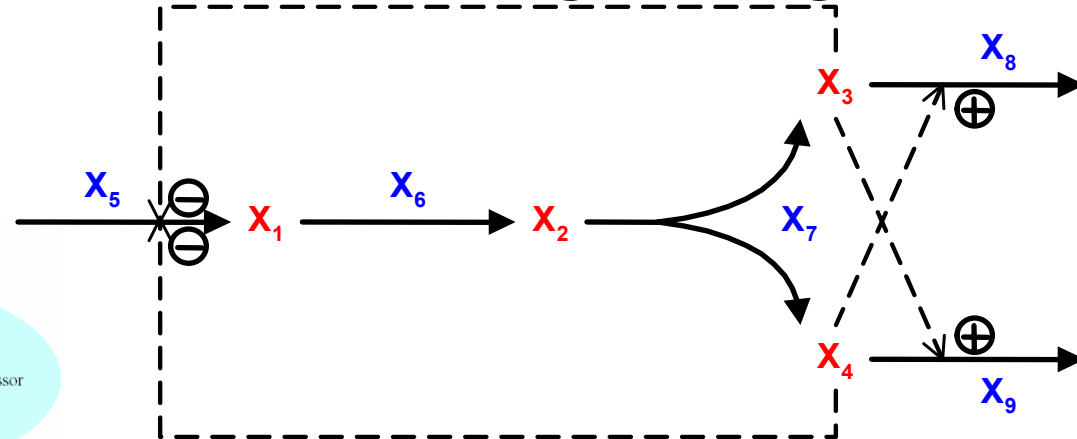


CHEMICAL- AND BIO-INFORMATICS

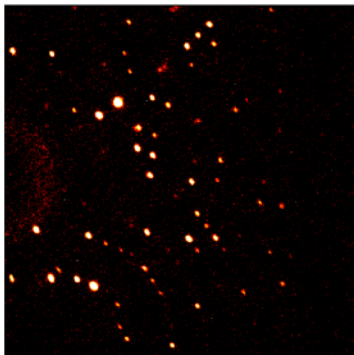
Molecular design



Metabolic engineering



X-ray imaging

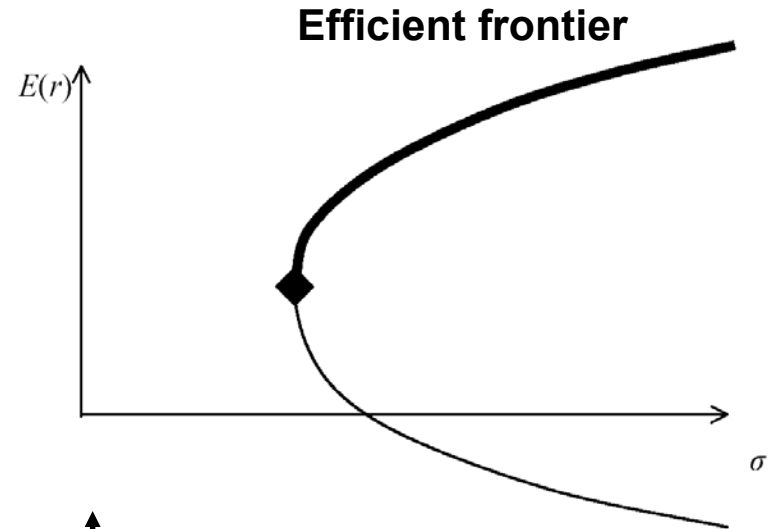
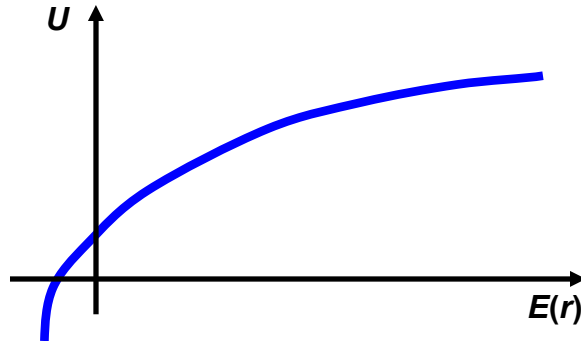


Bioinformatics

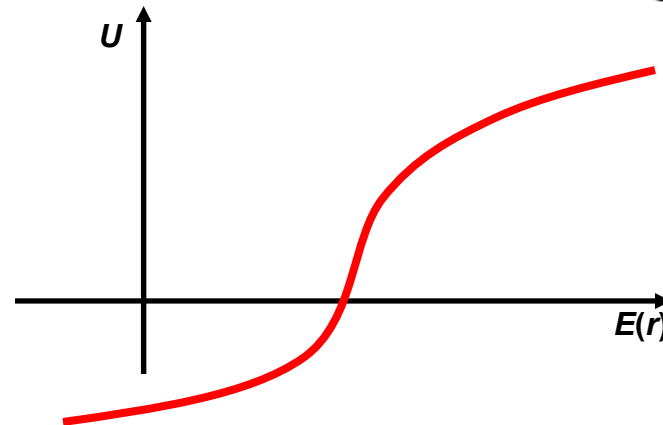
```
TTTGTGTCGTTTTCACAAAAATGGAAGTCCACA  
CACGGCGTCACACTTTGCTATGCCATAGCATTTT  
TACACAAAGTTAATAACGTGTCATGGTCATAT  
TGTATGCAAGGACGTCACATTACCGTCAGTAC  
GATCACGTTTTAGACCATTTTTTCGTCGTGAAC  
TTGTGATGTGTATCGATGTGTGTTGCCGAGTAGA  
TTTTCGCATCTTTGTTATGCTATGTTATTTTAT
```

PORTFOLIO OPTIMIZATION FOR WEALTH-DEPENDENT PREFERENCES

- **Markowitz model:**
 - variance vs. expected return
 - quadratic programming
 - risk-averse individuals



- **Empirical studies:**
 - Individuals have utility with two segments



NEW IN GAMS/BARON 6.0

- More decimals in solution, accurate timing of NLP sub-solvers, new branching and probing options, minimize and maximize, ...
- Provide dual solution
- Report lower bounds independent of tolerances
- Bug fixes
 - Infinite loops
 - » Local search
 - » `numsol > 1`
 - Power relaxations
 - x and $1/x$
 - *Thanks to:* Arne Drud, Oleg Shcherbina, Arnold Neumaier
- New options defaults
 - Reduced memory requirements
 - Faster convergence

COMPUTATIONAL RESULTS

- Collection of 102 NLP and MINLP problems from globallib and gamslib
- Average of 40.2 variables

	version 5	version 6	% Improvement
Nodes	2,187,512	76,112	96
Nodes to optimum	438,269	15,825	96
Nodes in memory	65,765	4,179	94
CPU sec	4,381	1,975	55

GAMS

www.gams.com

Global Optimization Solver BARON

Support

Sales

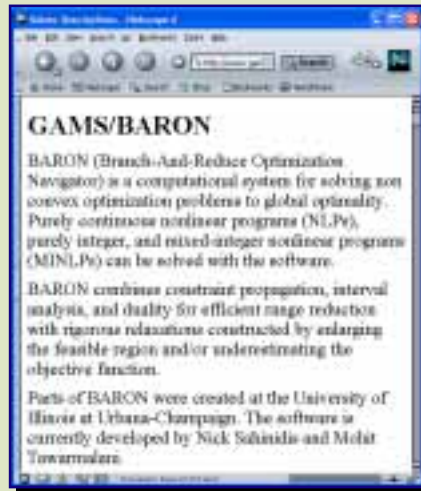
Solvers

Documentation

Model Library

Search

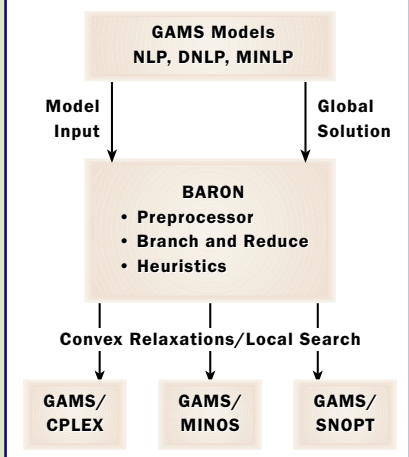
Contact Us



High performance and reliability of solvers come as a result of technological and theoretical developments in solution technology and modeling systems.

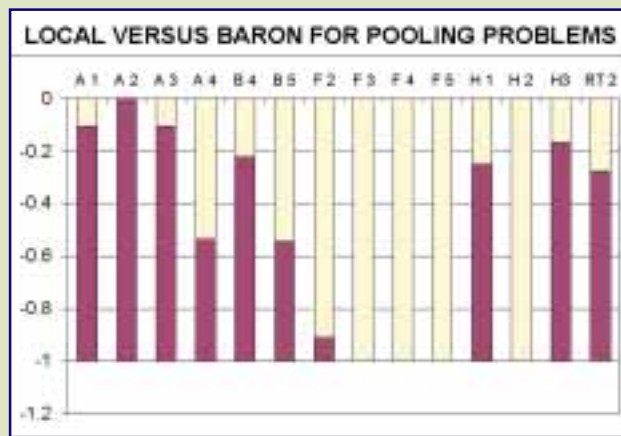
BARON combines constraint propagation, interval analysis, and duality for efficient range reduction while rigorous relaxations are constructed by enlarging the feasible region and/or underestimating the objective function.

Branch And Reduce Optimization Navigator



Important Application Areas for Global Optimization:

- Reliability
- Parameter Estimation
- Engineering Design
- Economics
- Finance
- Supply Chain Design and Operation
- Molecular Design and Bioinformatics



The darker area of each box is proportional to the difference between the global solution and the best local solution obtained using two local solvers for 14 pooling problems from the literature.

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