

Global Optimization and the GAMS Branch-and Cut Facility

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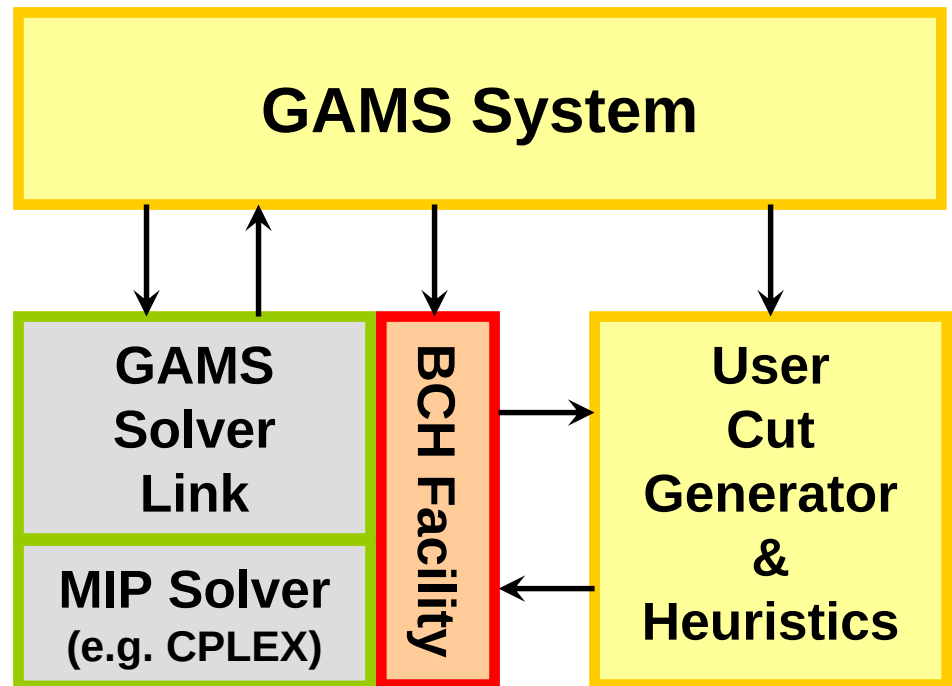
Branch-and-Cut (B&C)

- MIP problems benefit from *user supplied* cutting planes and integer solutions in a B&C algorithm
- Implementation facilities:
 - MIP solver callback functions (CPLEX, XPRESS, ...)
 - B&C framework (ABACUS, COIN BCP, ...)
- Required Knowledge for Implementation
 - IT knowledge (C/C++/JAVA, Solver APIs)
 - Mathematical programming knowledge
 - Application specific knowledge
- Supply GAMS users with an easy access to B&C:
GAMS Branch-and-Cut & Heuristic (BCH) Facility

BCH Unique Facility

- Cut Generator and Heuristic
 - Represented in terms of original GAMS problem formulation
 - Independent of the specific MIP solver
 - Use any other model type and solver available in GAMS

- Design Principle:



Example: Multi-Knapsack

<http://www.gams.com/modlib/libhtml/bchmknap.htm>

Binary variables $x(j)$; Positive variables $slack(i)$;
Equations $mk(i)$, $defobj$; Variable z ;

```
defobj.. z =e= sum(j, value(j)*x(j));
mk(i)..  sum(j, a(i,j)*x(j)) =l= size(i);
```

```
model m /all/; solve m max z using mip;
```

**The original
model formulation**

**Separation Problem
for Cover Cuts:**

$z.l < 1$

Cover Cuts $c(j)=y.l(j)$:
 $sum(c(j), x(j)) =l= card(j)-1$;

Binary variable $y(j)$ membership in the cover;
Equations $defcover$, $defobj$; Variable z ;

```
defobj..    z =e= sum(j, (1-x.l(j))*y(j));
defcover..  sum(j, ai(j)*y(j)) =g= size_i+1;
```

```
model cover /all/; solve cover min z using mip;
```

Cover Cut Separation Model

- Master model (bchmknapsack.gms):
`execute_unload 'multiknapsackdata', j, l, a, size // Export of data`
- Activate Cplex/BCH facility (cplex.opt):
`usercutcall bchcover.gms`
- Separation model (bchcover.gms):
`$gdxin multiknapsackdata // Get data from original model`
`$load j l a size`
`$gdxin bchout // Get fractional solution from solver`
`$load x`
`$include cutmodel // model definition and solution`
`* Define additional cut for original problem:`
`if (z.l<1, numcuts = 1;`
`x_c('1',j) = y.l(j); // cut matrix`
`rhs_c('1') = sum(j, y.l(j)) - 1;`
`sense_c('1') = 1); // 1 =l=, 2 =e=, 3 =g=`

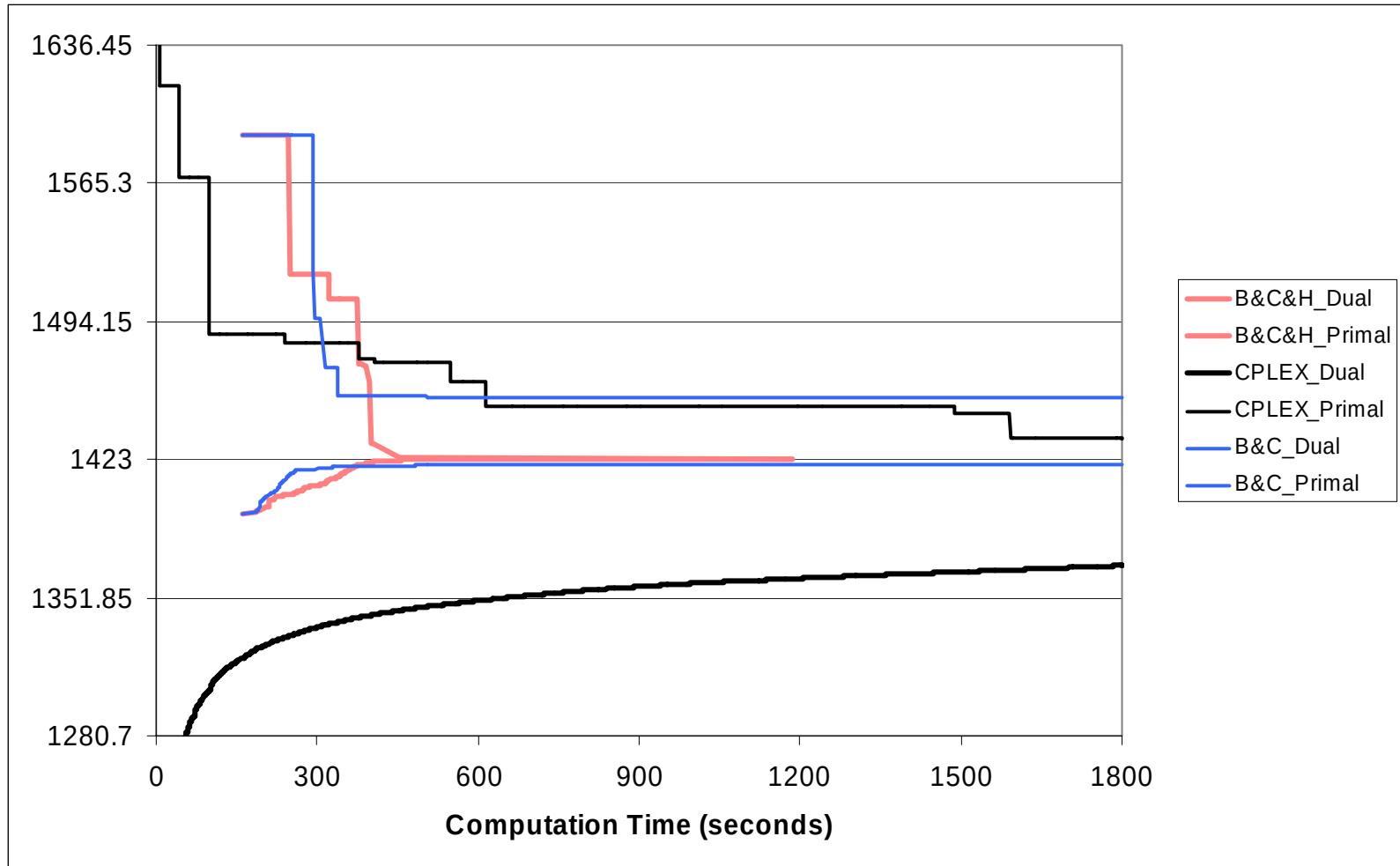
CPLEX with BCH Log

Nodes				Cuts/		ItCnt	Gap
Node	Left	Objective	IInf	Best Integer	Best Node		
0	0	4134.0741	2		4134.0741	3	
*** Calling cut generator. Added 1 cut							
		4090.6542	2		User: 1	4	
*** Calling cut generator. Added 1 cut							
		4001.4270	4		User: 1	5	
*** Calling cut generator. Added 1 cut							
		3950.4202	4		User: 1	6	
*** Calling cut generator. Added 1 cut							
		3871.4286	2		User: 1	7	
*** Calling cut generator. Added 1 cut							
		3841.6244	6		User: 1	9	
*** Calling cut generator. No cuts found							
*** Calling cut generator. No cuts found							
1	1	3800.0000	3		3800.0000	10	
*** Calling cut generator. Added 1 cut							
*** Calling cut generator. No cuts found							
					User: 1		
*** Calling cut generator. No cuts found							
*	2	1	0	3800.0000	3800.0000	12	0.00%
*** Calling cut generator. No cuts found							

Another Example

- Real Example: Oil Pipeline Design Problem
 - J. Brimberg, P. Hansen, K.-W. Lih, N. Mladenovic, M. Breton 2003. An Oil Pipeline Design Problem. Operations Research, Vol 51, No. 2 228-239
 - <http://www.gams.com/modlib/libhtml/bchoil.htm>
- Performance Improvements
 - Cplex/BCH: 20 minutes
 - Regular Cplex: 450 minutes
- Overhead of BCH
 - Time spent within the callback functions minus MIP computation on cuts and heuristics: 20% ~ 25%

Convergence – Pipeline Design



OK, but what has all this to do with Global Optimization?

- *Use any other model type and solver available in GAMS*
- Solve difficult MIP problems using GO for cut separation

Difficult Network Problem

- Single-commodity, uncapacitated, fixed-charge network flow problem

$$\min \sum_{(i,j) \in A} (f_{ij} y_{ij} + c_{ij} x_{ij}), s.t. \quad \sum_{(j,i) \in \delta^-(i)} x_{ji} - \sum_{(i,j) \in \delta^+(i)} x_{ij} = b_i, 0 \leq x_{ij} \leq M y_{ij}, y \in \{0,1\}$$

- Problem class includes Steiner Tree Problem
- F. Ortega, L. Wolsey A branch-and-cut algorithm for the single-commodity, uncapacitated, fixed-charge network flow problem. Networks 41 (2003), no. 3, 143—158
- 83 instances (including 24 not proven optimal)

Dicut Inequalities

- Dicut: $\sum_{(i,j) \in \delta^-(S)} y_{ij} \geq 1$ if $S \subset V$ and $b(S) > 0$.

- Separation problem:

$$\xi = \min \left\{ \sum_{(i,j) \in A} \bar{y}_{ij} z_j (1 - z_i) : \sum_{i \in V} b_i z_i > 0, z_i \in \{0, 1\} \forall i \in V \right\}$$

- Non-convex quadratic binary program:
 - Ortega/Wolsey: Greedy algorithm
 - Let's solve this exactly (small, #variables $|V|$)

Global versus Local

- Berlin52 – from SteinLib
 - 52 nodes (1 source, $n'=15$ sinks), 1326 edges

	CPLEX (no BCH)	BCH + SBB (local)	BCH + BARON (global)
# cuts	0	139	610
# nodes	168449	1	1
Time	1000	1000	223
Gap	45%	39%	0%

Heuristics

- Local Branching (Fischetti/Lodi) or Delta/Onion Method (Guignard/Spielberg)
- Improve a given binary solution x^* by restricting search space. Add inequality:

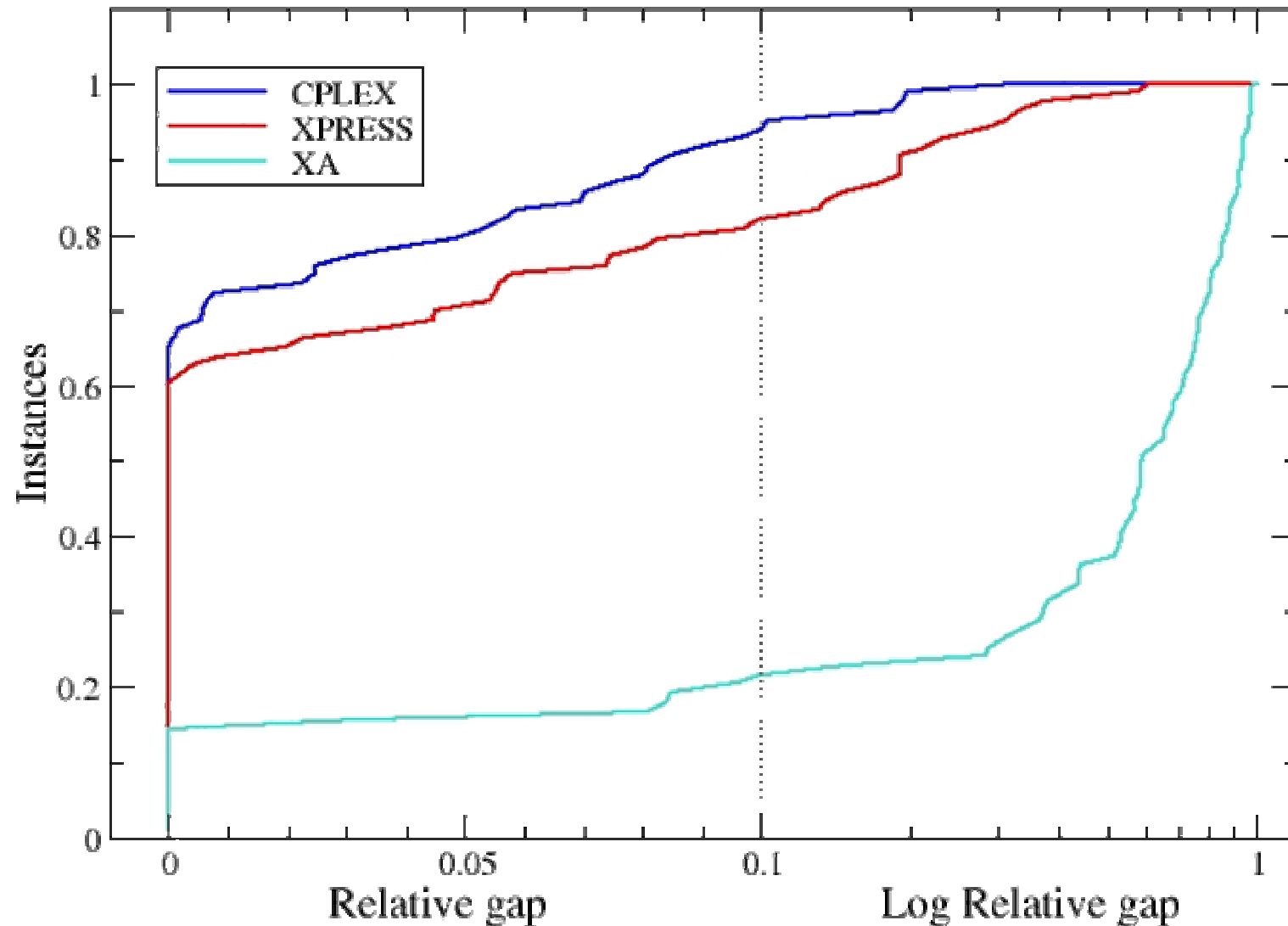
$$\sum_{j, x_j^*=1} x_j + \sum_{j, x_j^*=0} (1 - x_j) \leq k$$

- Solve resulting problem as regular MIP
- Solve recursively/iteratively

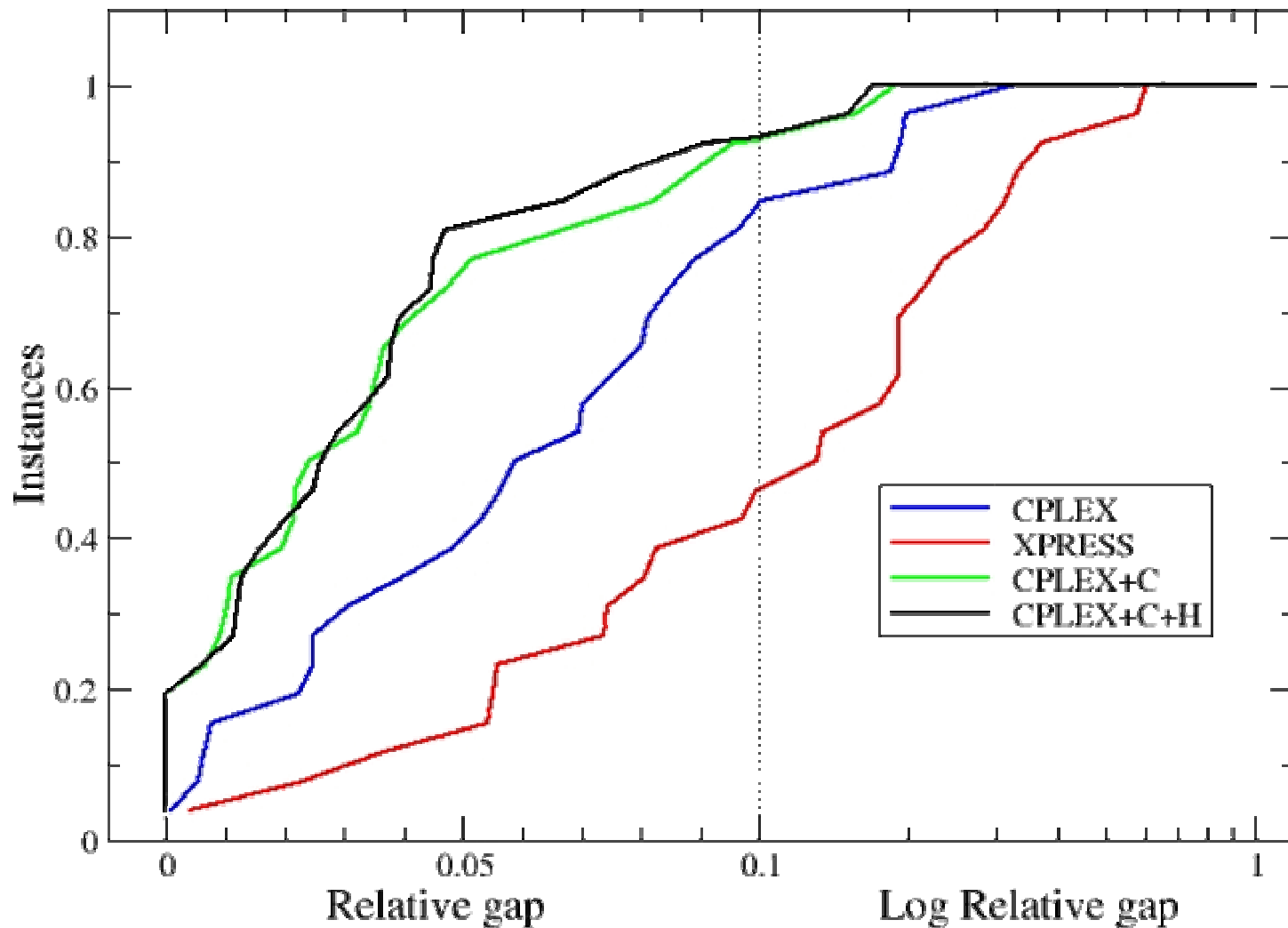
Computational Experiments

- 83 instances from Wolsey's web page:
<http://www.core.ucl.ac.be/wolsey/ufcn.htm>
- Models in Xpress-Modeler & MPS format
- Translated in consistent GAMS model with 83 data files
- Verified translated instances:
 - Same relaxed objective value
 - Matrix statistics
 - When solved optimally, same objective value
- Computing Environment
 - Athlon 2800, 2 GB memory running Linux
 - GAMS 21.4, Cplex 9.0, Xpress 2003
 - Time limit 1800 seconds

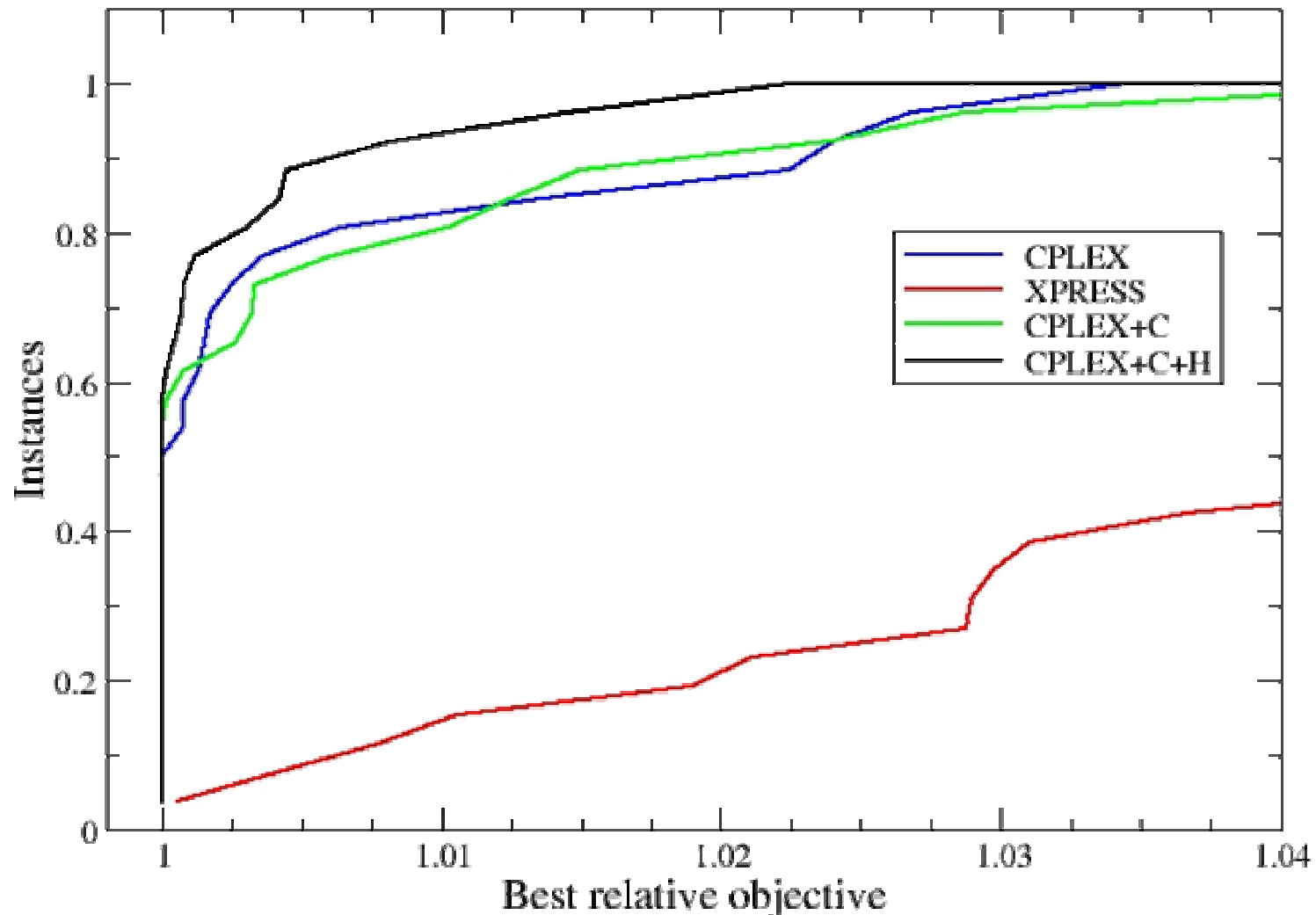
No BCH – 1800 s Time Limit



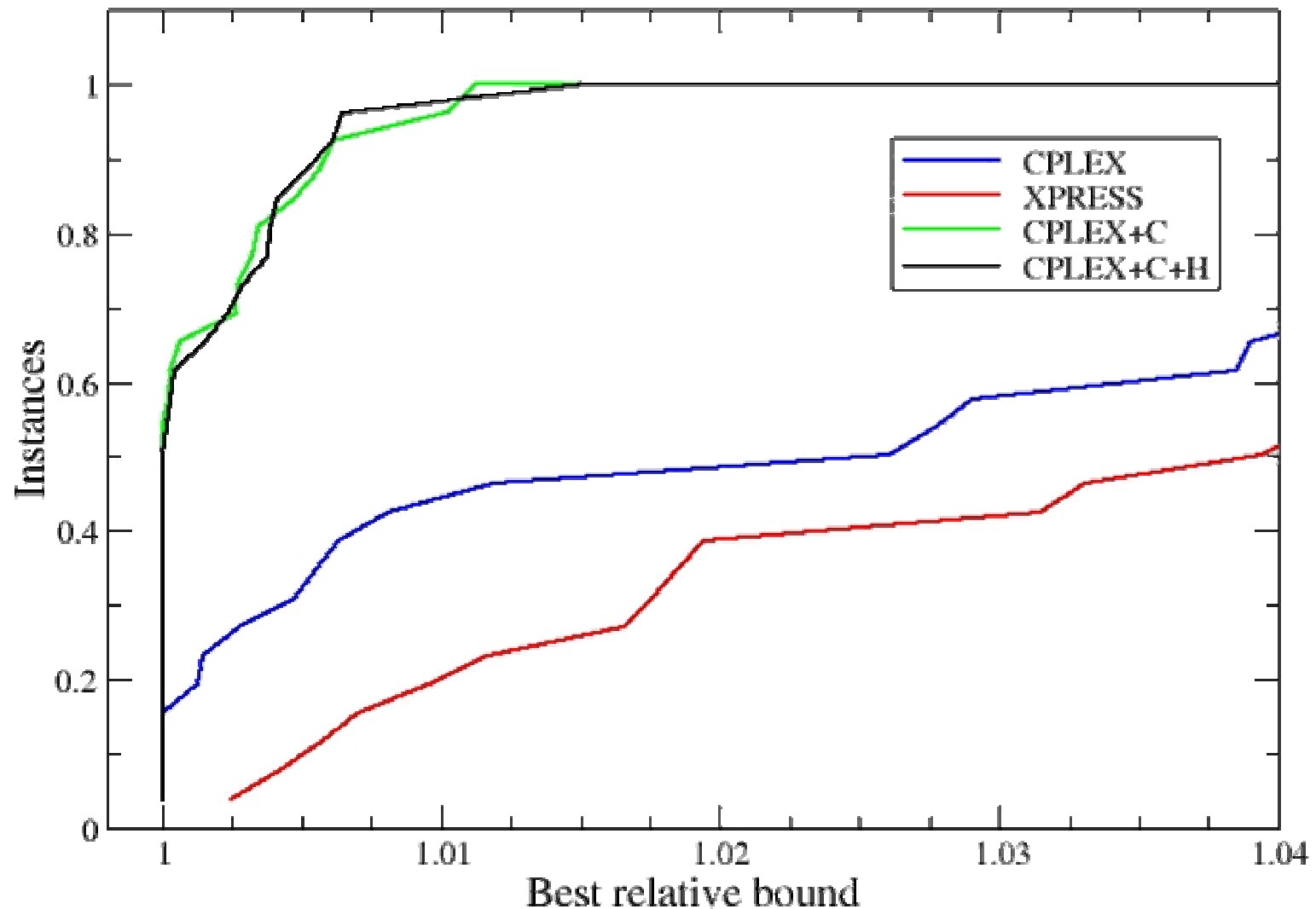
Unsolved Instances - Gap



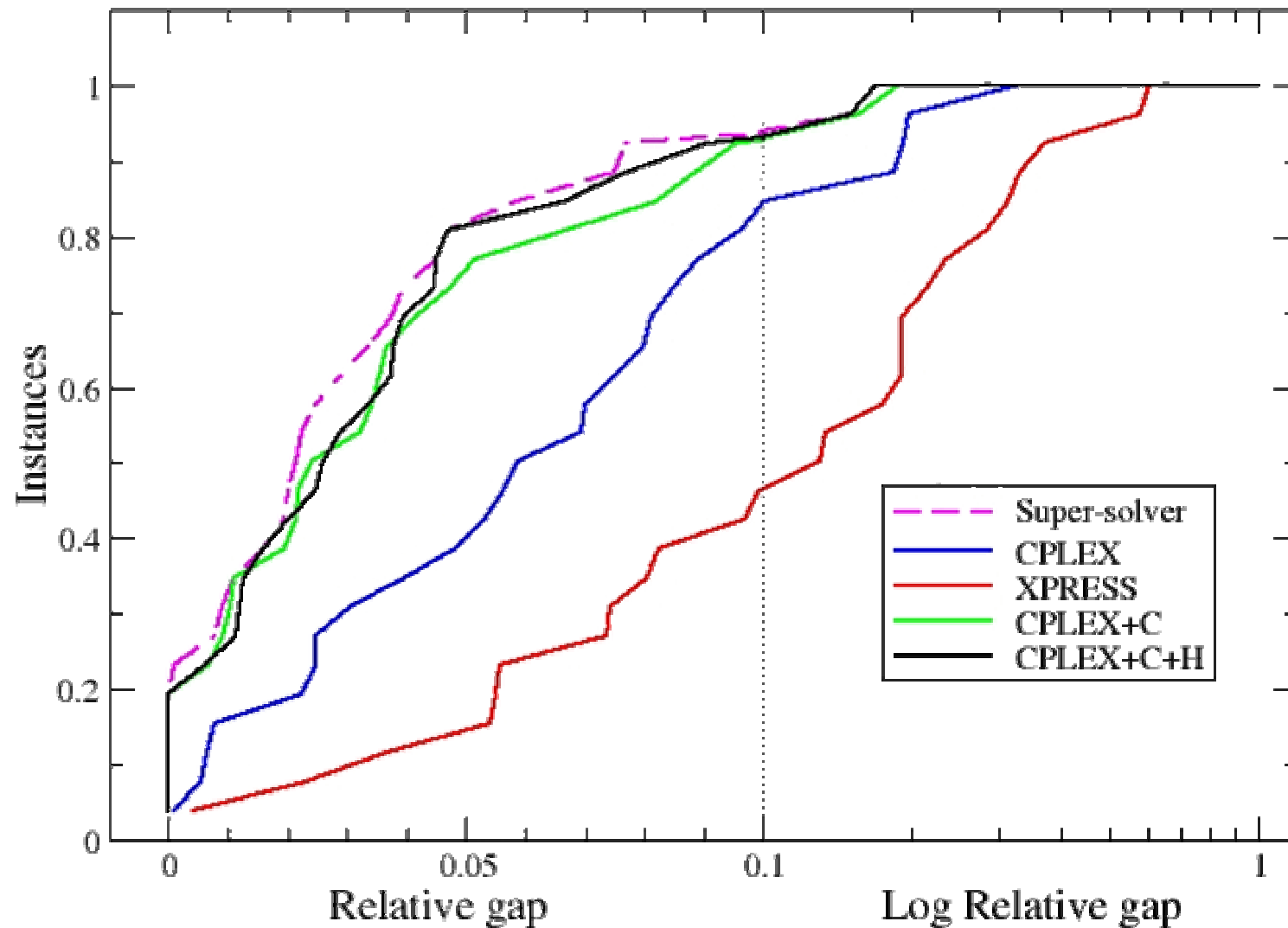
Unsolved Instances - Objective



Unsolved Instances - Bound



Unsolved Instances - Gap



Conclusion

- BCH Facility automates all major steps necessary to define, execute, and control the use of user defined routines within the framework of general purpose MIP codes
- Cut separation models provides a new application area for Global Optimization
- More information:
 - www.gams.com/docs/bch.htm (BCH)
 - www.gams.com/apps/fcnet.htm (FC-Network)
 - www.gams.com/presentations (Presentation)